

17/6/21

Agenda: Assignment 7 questions & answers
 • Some general talk around mean & variance.

(Q1) • $E[X] = \sum_{n=1}^{\infty} n P(X=n)$.
 • $Var(X) = E[X^2] - E[X]^2$.

(a) $\Omega = \{1, 2, \dots, n\}$, $P(\omega) = 1/n \forall \omega$,
 $X(i) = i, i \in \Omega$
 $\Rightarrow E[X] = \sum_{i=1}^n i/n = \frac{n(n+1)}{2n} = \frac{n+1}{2}$.
 $E[X^2] = \sum_{i=1}^n i^2/n = \frac{n(n+1)(2n+1)}{6n} = \frac{(n+1)(2n+1)}{6}$.
 $\Rightarrow Var(X) = \frac{(n+1)(2n+1)}{6} - \left(\frac{n+1}{2}\right)^2 = (n+1) \left[\frac{2n+1}{6} - \frac{n+1}{4} \right]$
 $= (n+1) \left[\frac{n-1}{12} \right] = \left[\frac{n^2-1}{12} \right]$

(b) $\Omega = \{a_1, a_2, \dots, a_n\} = \{a_i \in \{0, 1\}^n : P(\omega) = 1/2^n \forall \omega \in \Omega$
 $X(\omega) = \# \{i : a_i = 1\} \Rightarrow P(X=k) = \binom{n}{k} 2^{-n}, 0 \leq k \leq n$
 Let $X_i(\omega) = 1_{\{a_i=1\}} \Rightarrow X_i$ are independent and
 $X = \sum_{i=1}^n X_i$.
 $\Rightarrow Var(X) = \sum_{i=1}^n Var(X_i)$.
 $E[X_i] = P(X_i=1) = p; E[X_i^2] = P[X_i=1] = p$
 $\Rightarrow Var[X_i] = P(1-p)$.
 $\Rightarrow Var(X) = np(1-p)$.

(c) Description of experiment:
 - Collection of n objects
 - Of which m are good.
 - Draw n at a time, and count the no. of successes.
 To write this cleanly, let $1, \dots, m$ be "good" and $m+1, \dots, n$ be "bad". Then, we have the sample space
 $\Omega = \{S \subseteq [n] : |S|=n\}; P(S) = 1/\binom{n}{n}$.
 $X(S) = \# \{i \in S : i \leq m\}$.
 Then, $P(X=t) = \frac{\binom{m}{t} \binom{n-m}{n-t}}{\binom{n}{n}}$
 $\forall \max\{0, n-m\} \leq t \leq \min\{m, n\}$

Idea: For $k \leq m, X_i = 1_{\{i \in S\}}$. Then, $\sum_{i=1}^m X_i = X$, note that X_i need not be ind. b/c it is $i-1$ sampling without replacement.

However, $E[X_i] = P(i \in S) = \frac{m}{n}$; so $E[X] = \frac{m \cdot n}{n}$.

$E[X_i^2] = P(i \in S) = \frac{m}{n}$.

$Var(X) = E[X^2] - E[X]^2 = E[\sum_{i=1}^m X_i^2] - (E[X])^2$

$= \sum_{i=1}^m E[X_i^2] - \sum_{i,j} E[X_i] E[X_j]$

$= \sum_{i=1}^m (E[X_i^2] - E[X_i]^2) + \left(\sum_{i,j} E[X_i X_j] - E[X_i] E[X_j] \right)$

$= \frac{m \cdot n}{n} \left(1 - \frac{m}{n} \right) m + \dots$

Now, for $i \neq j$, $E[X_i X_j] = P(i, j \in S) = \frac{\binom{n-2}{n-2}}{\binom{n}{n}} = \frac{n}{n} \times \frac{n-1}{n-1}$.

The number of $i \neq j$ is $\frac{m(m-1)}{2}$ so we get:

$= \frac{m \cdot n}{n} \left[1 - \frac{m}{n} \right] m + \frac{m(m-1)}{2} \left[\frac{2 \cdot n(n-1)}{n(n-1)} - \frac{n^2}{n^2} \right]$.

Can be simplified if necessary.

(Q2) $\Omega = [6] \times [6], P(\omega) = 1/36 \forall \omega \in \Omega$.

(a) $X(a,b) = a+b$

Expand in detail to notice

$P(X=2) = P(X=12) = 1/36, P(X=3) = P(X=11) = 2/36,$

$P(X=4) = P(X=10) = 3/36, P(X=5) = P(X=9) = 4/36,$

$P(X=6) = P(X=8) = 5/36, P(X=7) = 5/36$

i.e. $P(X=a) = \frac{6-|7-a|}{36}, a \in \{2, 3, \dots, 12\}$.

Note that $X = X_1 + X_2$, where $X_1(a,b) = a$

$X_2(a,b) = b$

$\Delta X_1, X_2$ are ind. $\Rightarrow E[X] = 3.5 + 3.5 = 7$.

Note $X_1, X_2 \sim \text{Unif}[6]$

$\Rightarrow Var(X) = 2 \left(\frac{6^2-1}{12} \right) = 35/6$.

(b) $X(a,b) = \max\{a, b\}$.

$\Rightarrow P(X=k) = \frac{2k-1}{36}, k=1, 2, \dots, 6$.

$E[X] = \sum_{k=1}^6 \frac{k(2k-1)}{36} = \frac{1}{36} \sum_{k=1}^6 (2k^2 - k)$

$= \frac{1}{36} \left[\frac{6 \cdot 7 \cdot 13}{3} - 21 \right] = \frac{1}{36} [161] = 161/36 \approx 4 \frac{17}{36}$

$E[X^2] = \sum_{k=1}^6 \frac{k^2(2k-1)}{36} = \frac{1}{36} \left[\sum_{k=1}^6 (2k^3 - k^2) \right]$

$= \frac{1}{36} \left[\frac{36 \cdot 49}{2} - 21 \right] = \frac{1}{36} [861] = 861/36$.

$\Rightarrow Var(X) = 861/36 - \left(161/36 \right)^2 = 5075/1296$.

(Q5) $\Omega = \{x_1, \dots, x_n\} : x_i \geq 0, \sum_{i=1}^n x_i = n$.

$P(x_1, \dots, x_n) = \frac{1}{n!} \frac{n!}{x_1! x_2! \dots x_n!}$

$\rightarrow X(x_1, \dots, x_n) = \# \{i : x_i = 0\}$.

Write $X_i = 1_{\{x_i=0\}}$, then $X = \sum_{i=1}^n X_i$.

Now, we will find $P(X_1=1, \dots, X_j=1) = P(x_1=0, \dots, x_j=0)$

$= \sum_{\substack{x_1=0, \dots, x_n=n \\ x_i=0 \forall i}} \frac{1}{n!} \frac{n!}{x_1! \dots x_n!} = \sum_{i=0 \dots n-j} \frac{1}{n!} \frac{n!}{i! \dots (n-j)!}$

$= \frac{1}{n!} (n-j)^n$ by the multinomial thm.

Now use the generalized IEP:

$\{X=k\}_n = \{ \text{For exactly } k \text{ values of } i, X_i=1 \}$

$= \sum_{j=k}^n (-1)^{j-k} \binom{j}{k} \binom{n}{j} \left(\frac{n-j}{n} \right)^n$.

This gives us the pmf of X . However, we can get the expectation from:

$E[X] = \sum E[X_i] = \sum P(X_i=1) = n \left(\frac{n-1}{n} \right)$

And the variance from:

$E[X^2] = \sum E[X_i^2] + \sum_{i \neq j} E[X_i X_j]$

$= \sum P[X_i=1] + \sum_{i \neq j} P[X_i=1, X_j=1]$

$= n \left(\frac{n-1}{n} \right) + 2 \binom{n}{2} \left(\frac{n-2}{n} \right)^n$

and finish.

(Q7) This is fair! Com-ntable: Let X take distinct values $\{a_1, a_2, \dots, a_n\}; a_i \geq 0$. Let

t be positive. Let $a_0 = 0, a_{n+1} = \infty$ (formally). Then, for t positive, we know

$a_i \leq t \leq a_{i+1}$ for some $i \in \{0, \dots, n\}$. So, we get:

$t P(X \geq t) \leq \sum_{j=i+1}^n P(X=a_{j+1}) a_{j+1} \leq E[X]$

from where Markov follows. Chebyshev follows from:

$P(|X-E[X]| \geq t) = P[(X-E[X])^2 \geq t^2] \leq \frac{E[(X-E[X])^2]}{t^2}$

$= Var(X)/t^2$

and done!

A look at variance - mean comparison.

The strength of Chebyshev's inequality is measured by the amount of variance of X . Ideally, X is expected to concentrate around its mean, if the variance is of order smaller than the mean.

Two examples

$\rightarrow U(n) \rightarrow E[X] \approx n, Var[X] \approx n^2$.

$\rightarrow \text{Bin}(n, p) \rightarrow E[X] \approx n, Var(X) \approx n$.

If you look at the Chebyshev inequality,

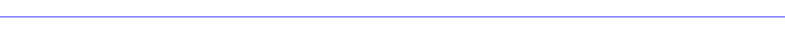
$P(|X-E[X]| \geq t) \leq Var(X)/t^2$.

Let $t = C \sqrt{Var(X)}$, then

$P(|X-E[X]| \geq C \sqrt{Var(X)}) \leq 1/C^2$.

Now, the smaller the quantity $Var(X)$ in comparison to the expectation, the better the applicability of Chebyshev. So, we expect Binomial to concentrate better than a uniform.

Indeed, from the pmf's,

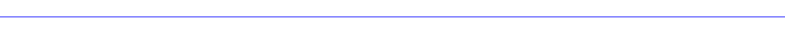


it is clear that the binomial is much closer to its mean than that of a uniform.

The point: knowing just the exp. & variance gives a great idea of the concentration of the r.v., without even knowing the pmf.

Briefly Monte-Carlo

Imagine you have a shape whose area you want to find. Let's say it's a subset of a unit square.



The idea is to have a shooter, whose job is to shoot bullets at the target, in a manner such that they are uniformly spread over Ω . Here's a question: If N is the number of shots, & n is the no. of shots that land inside S . Is $\frac{n}{N} \approx \text{Area}(S)$?

Ans: Yes. That is the content of the laws of large numbers, and forms the basis of the Monte Carlo method.

Qn: How large an N do I need, to get a good idea of the area of S ?

Ans: Depends upon the Variance of the uniform random variable.

In general, a "Monte Carlo" scheme involves calculating a deterministic quantity, using independent identically distributed random variables. The convergence of the Monte-Carlo scheme involves the mean-variance relation of the random variable.

So, a MC scheme involving a Binomial r.v. is likely to converge faster & be accurate quicker than one involving a uniform r.v.

The idea of the mean and variance

Suppose you have points scattered on a plane, at positions v_i with mass m_i . Think of this as a "rigid" mass.

Qn: Imagine I had to interpret this as a "point" mass. Which point is the best for this purpose?

Answer: You try to minimize a "reasonable" distance function like so:

$f(x) = \sum m_i |x - v_i|^2$

This is the "squared distance" function.

Fact: If X is a random variable with $P(X=v_i) = m_i$, then $E[X]$ is the unique minimizer of f . In other words, $E[X]$ is the best place for the whole rigid mass to be placed as a point mass. This is the "geometric" interpretation of $E[X]$.

Translation: Without knowing anything about a random variable, the best guess you can take about it's values, is it's expectation.

If we knew more, can we do better though? For functions $f, g: \{1, \dots, n\} \rightarrow \mathbb{R}$, define the "squared distance"

$\|f-g\| = \sum (f(x) - g(x))^2$

If we treat a function as a point, by taking $f(x) = (f(1), \dots, f(n)) \in \mathbb{R} \times \mathbb{R} \times \dots \times \mathbb{R}$, n times,

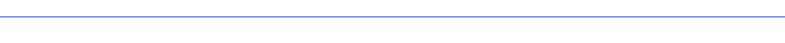
then the distance between functions is given a geometric meaning as the distance between 2 points in a high dimensional space.

So pmf's, are also "points" in a high dimensional space. With this, "geometrically" speaking:

If X is a random variable with pmf f , then conditioning X w.r.t. some other event, introduces a "set" of possible pmf's that can be the conditional pmf. Now, the actual cond. pmf. will be given by:

$f_c(x) = \argmin_{g \in S} \|f-g\|$

That is geometrically, the closest point in S , to f , is the conditional pmf!



Further: The distances, in each case, are given by the variance & conditional variance. So, in a sense:

— Expectation = Best approx., Variance = Error of Approx.

— Condnl. Expn = Best approx. when the pmf's are restricted to a certain set, & Condnl. Variance = Error of this Approx.